

# Quantum Precision: Unbounded memory advantage in stochastic simulation using quantum mechanics.

Andrew J. P. Garner  
*10<sup>th</sup> January 2017*

*Workshop on the interdisciplinary frontiers  
of quantum and complexity science 2017*

arXiv:1609.04408



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# Collaborators



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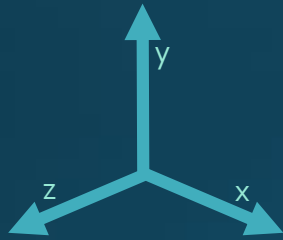


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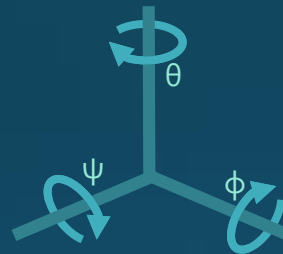
# Motivation

# Real numbers are ubiquitous

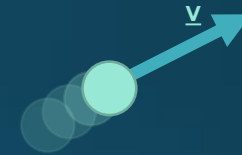
We often want to simulate systems whose properties are described by real numbers:



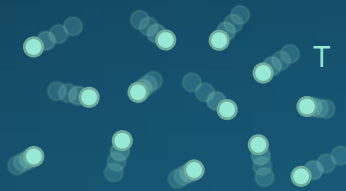
Positions



Angles



Velocities



Temperatures

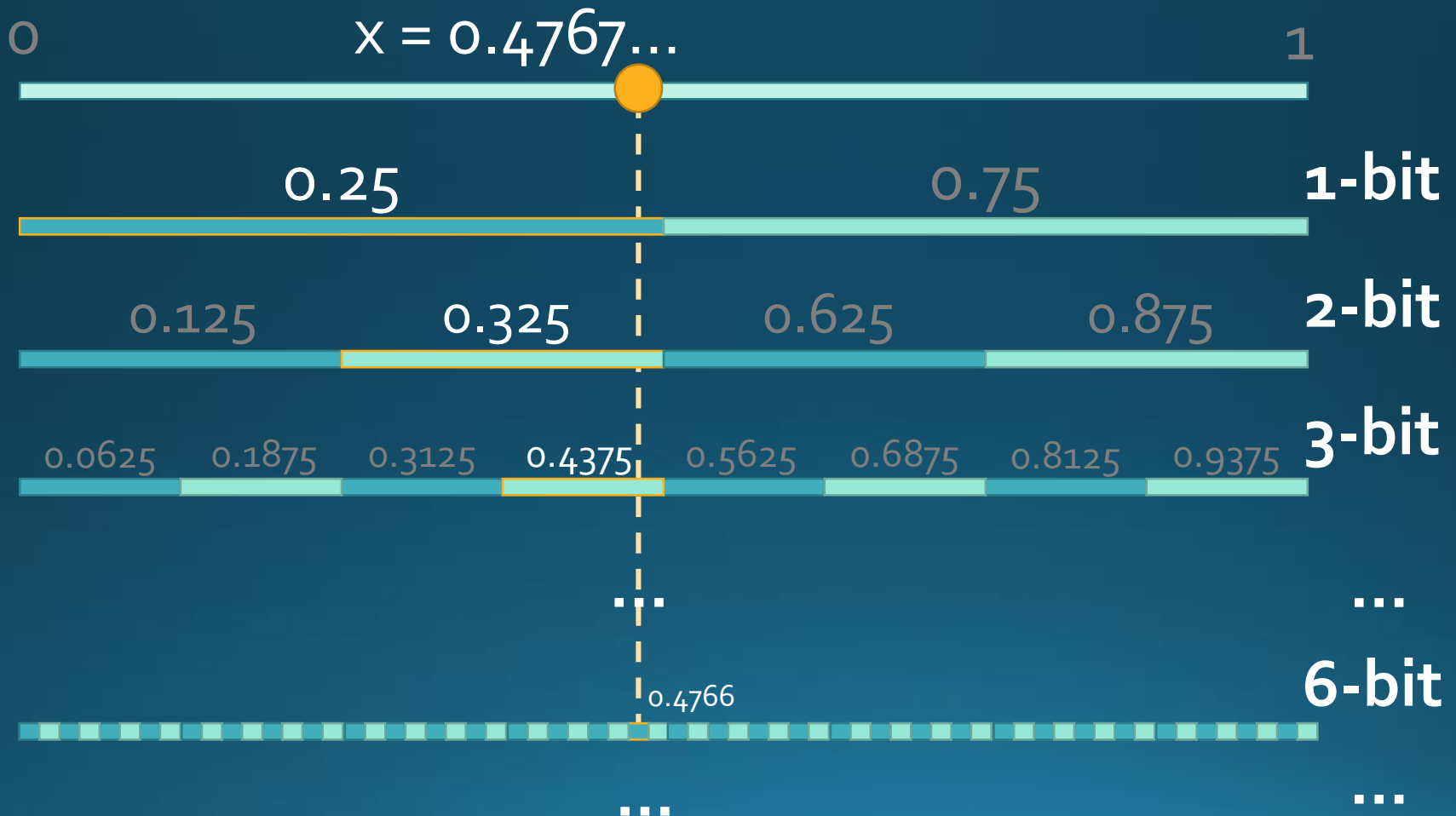


Field strengths

# Digital computers cannot directly store real numbers



# Precision is memory-limited



# Precision and memory cost

There are two subtly distinct concepts:

**Precision** – the closeness of a number to some ‘true’ value – quantified by the number of divisions  $N = 2^n$  (statistical property of simulation)

**Memory cost** – the amount of information required to store a value (physical property of simulator)

In classical information theory, these are so fundamentally related that they are used almost interchangeably (e.g. “32-bit precision”)

In quantum mechanics these two quantities can diverge significantly!

# Scenario



# Scenario: cyclic random walks

Consider a bead on a ring:

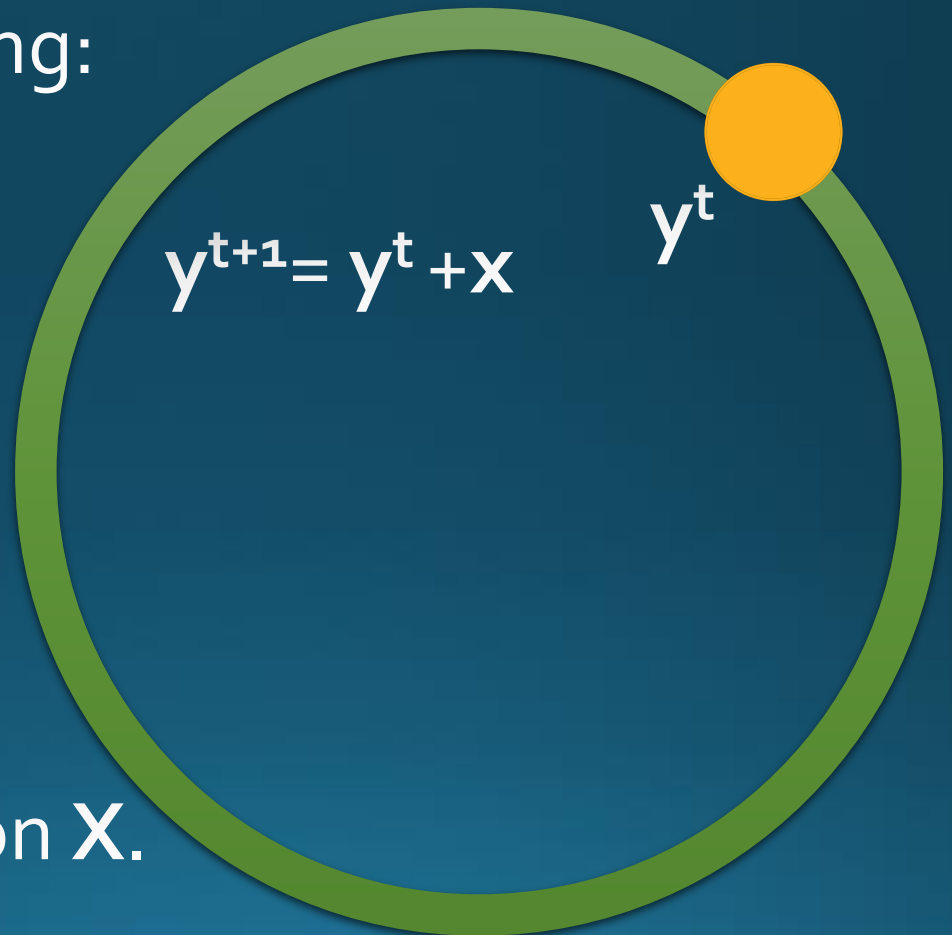
Position  $\mathbf{y}^t$  at time  $\mathbf{t}$ .

Discrete time-steps.

Position at time  $\mathbf{t}+1$ :

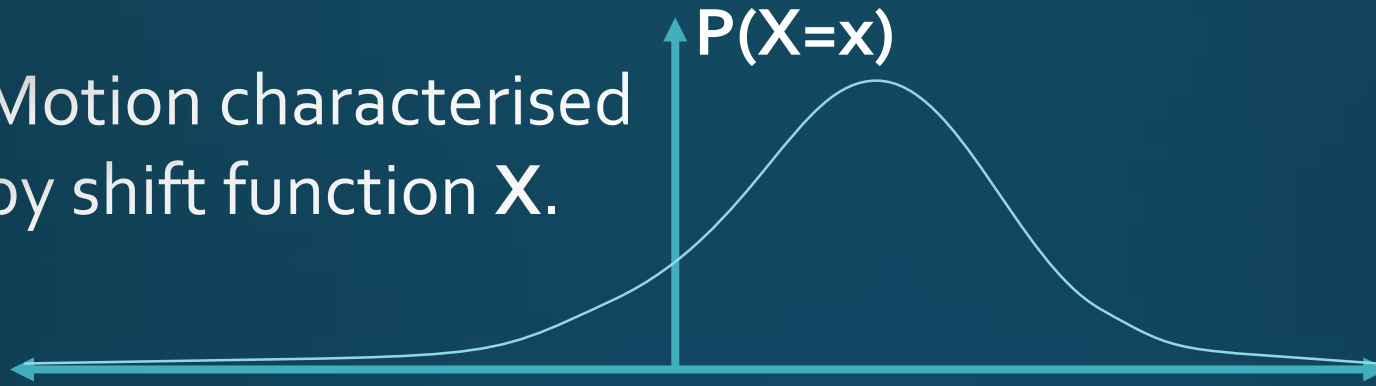
$$\mathbf{y}^{t+1} = \mathbf{y}^t + \mathbf{x}$$

where  $\mathbf{x}$  is randomly chosen from some (continuous) distribution  $\mathbf{X}$ .

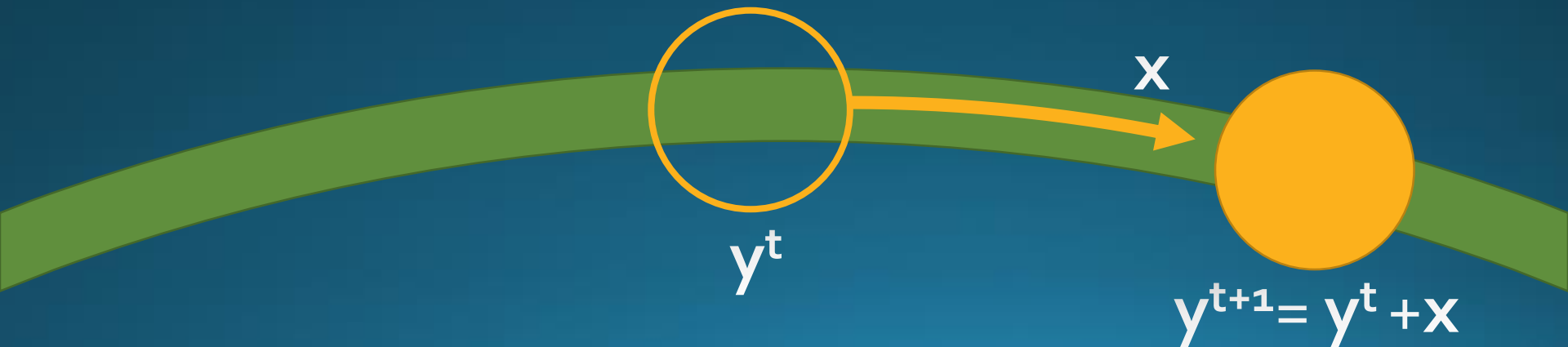


# Scenario: cyclic random walks

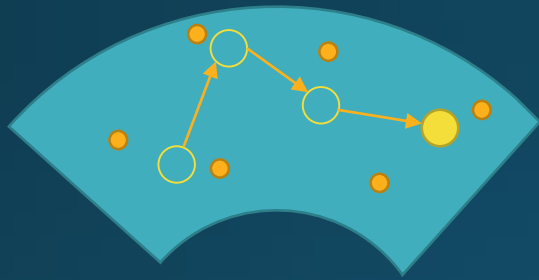
Motion characterised  
by shift function  $X$ .



Path described by series of continuous random  
variables  $Y^0, Y^1 \dots Y^t \dots$  obeying  $Y^{t+1} = Y^t + X \pmod{1}$



# Scenario: cyclic random walks



Azimuthal  
motion of  
diffusing gas



Bead  
on ring



Single  
electron in  
circuit

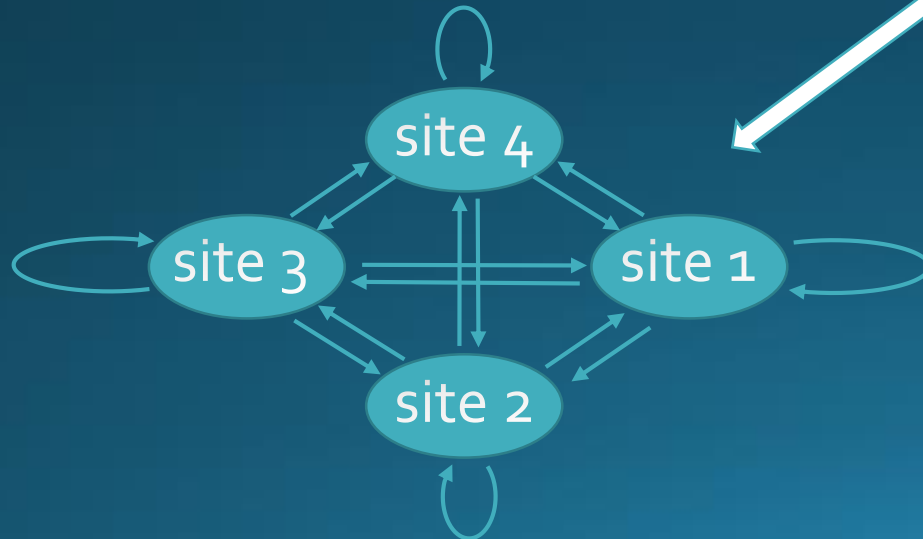
Any system with symmetry and diffusion

# Classical simulators

To simulate to precision  $n$ , the ring is split into  $N=2^n$  sites,  $\{s\}$ .

Shift function  $X$  approximated as  $N \times N$  *stochastic matrix*.

Simulation is effectively a finite state machine:



Statistics of future positions  $\gamma^{t+1}, \gamma^{t+2}, \dots$  depend on the current site  $s^t$ .

This information must be stored in the simulator's memory!

# The cost of classical simulation

More divisions gives greater statistical accuracy of simulation.

The state machine's memory must accommodate more possibilities.

For process with randomness, in steady state each site is equally likely.

*Entropy cost* of simulating to precision  $n = \log_2 N$  is then:

$$H_{\text{classical}} = n \quad (\text{bits})$$



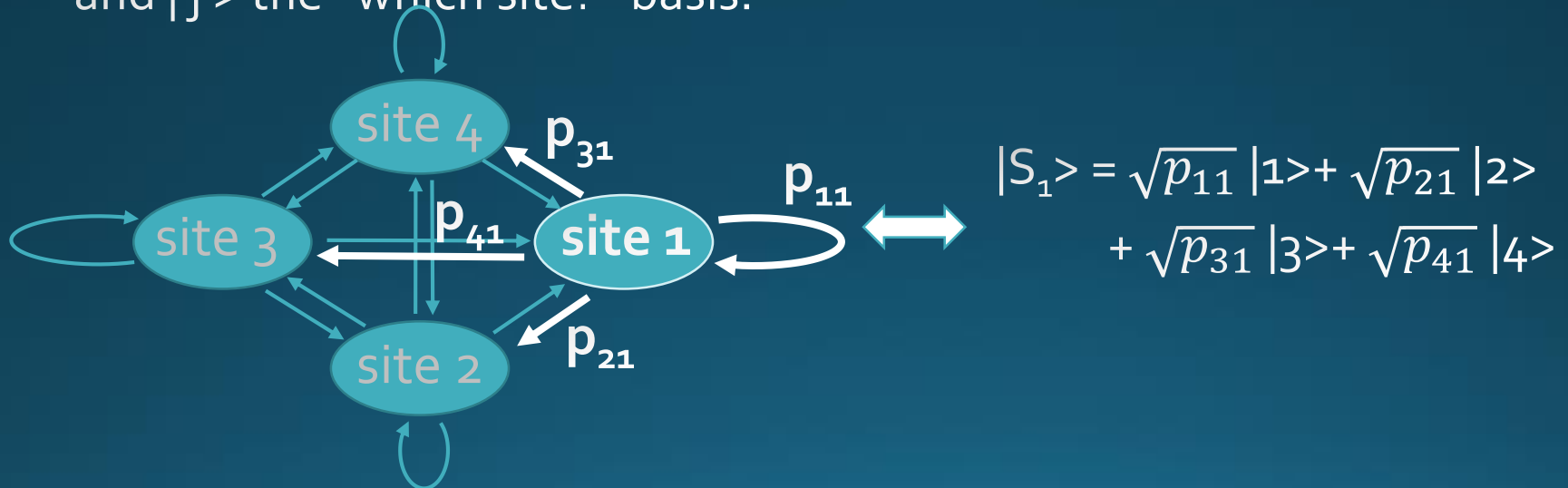
# Quantum simulator

# Quantum simulator

For a given precision  $n$ , each site is mapped onto a quantum state:

$$|S_i\rangle = \sum_{j=1}^{2^n} \sqrt{p_{ji}} |j\rangle$$

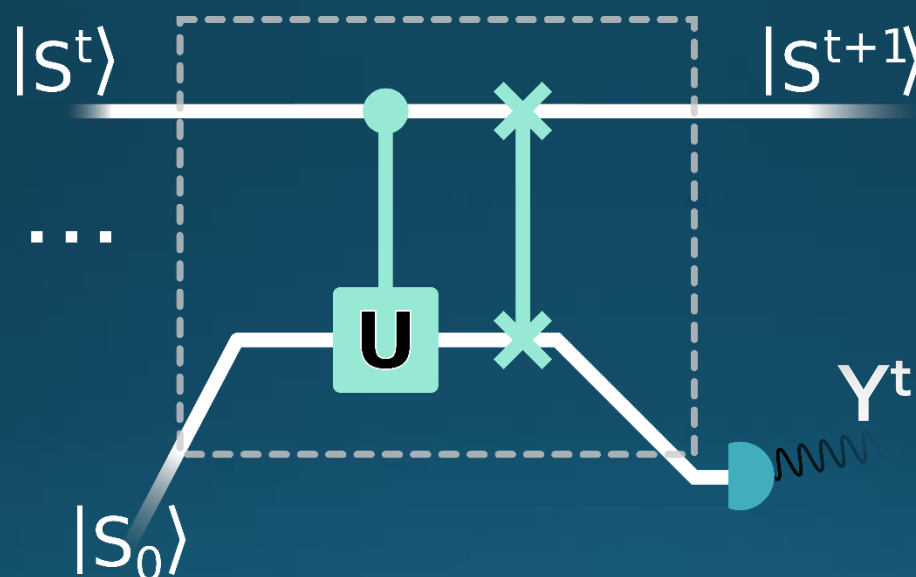
where  $p_{ji}$  is the probability of transitioning from site  $i$  to  $j$ , and  $|j\rangle$  the “which site?” basis.



Measurement of  $|S_1\rangle$  in the ‘which site?’ basis collapses state onto one of the four possible sites with the correct probability.

# Quantum simulator

When supplied with appropriate state  $|S^t\rangle$ , the statistics for  $Y^t$ , and the next state may be generated by the following circuit:

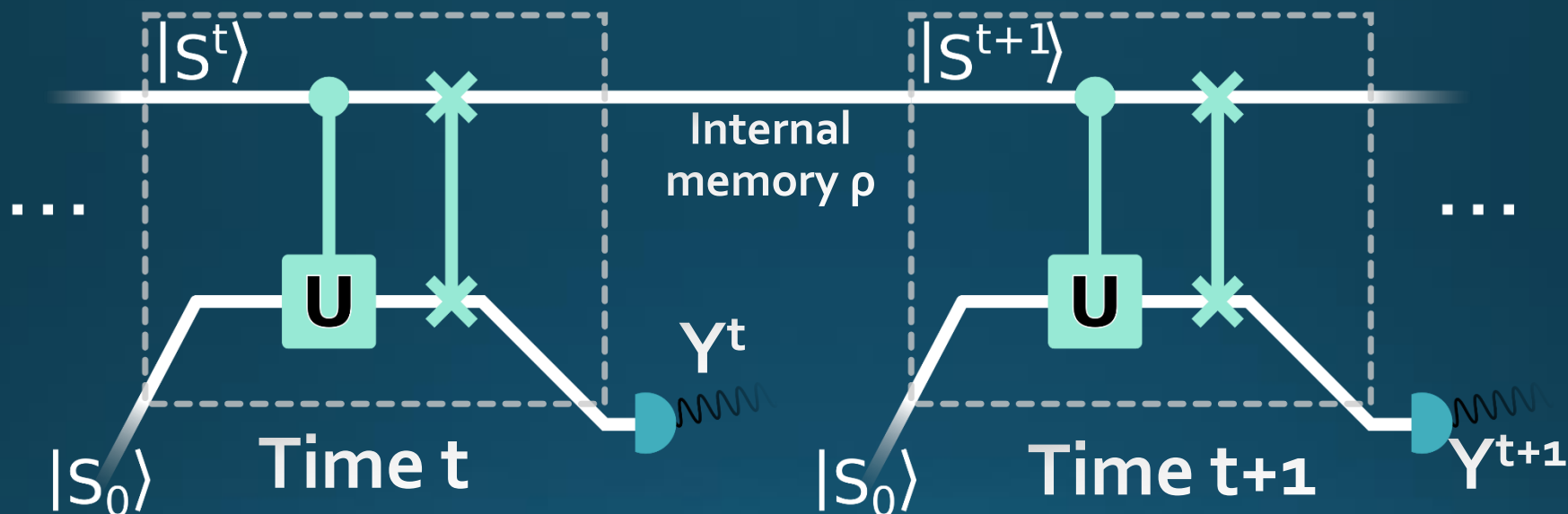


where  $U$  is a controlled unitary that for each  $i$  rotates initial state  $|S_0\rangle$  to state  $|S_i\rangle$  when the control bit is in state  $|i\rangle$ .



# Quantum simulator

Internal memory  $\rho$  persists between time-steps:



Steady-state of internal memory is  $\rho = \frac{1}{2^n} \sum_{i=1}^{2^n} |S_i\rangle \langle S_i|$ .

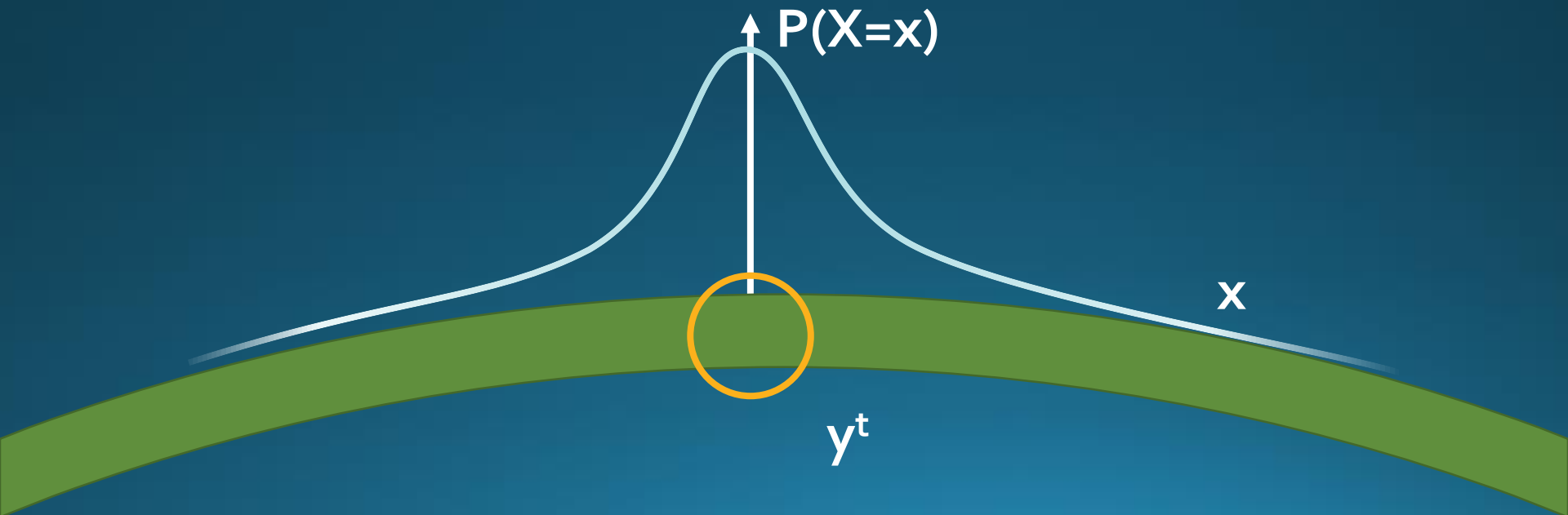
Memory requirement quantified by *von Neumann* entropy:

$$H = -\text{Tr} \rho \log \rho$$

# Example: Gaussian noise

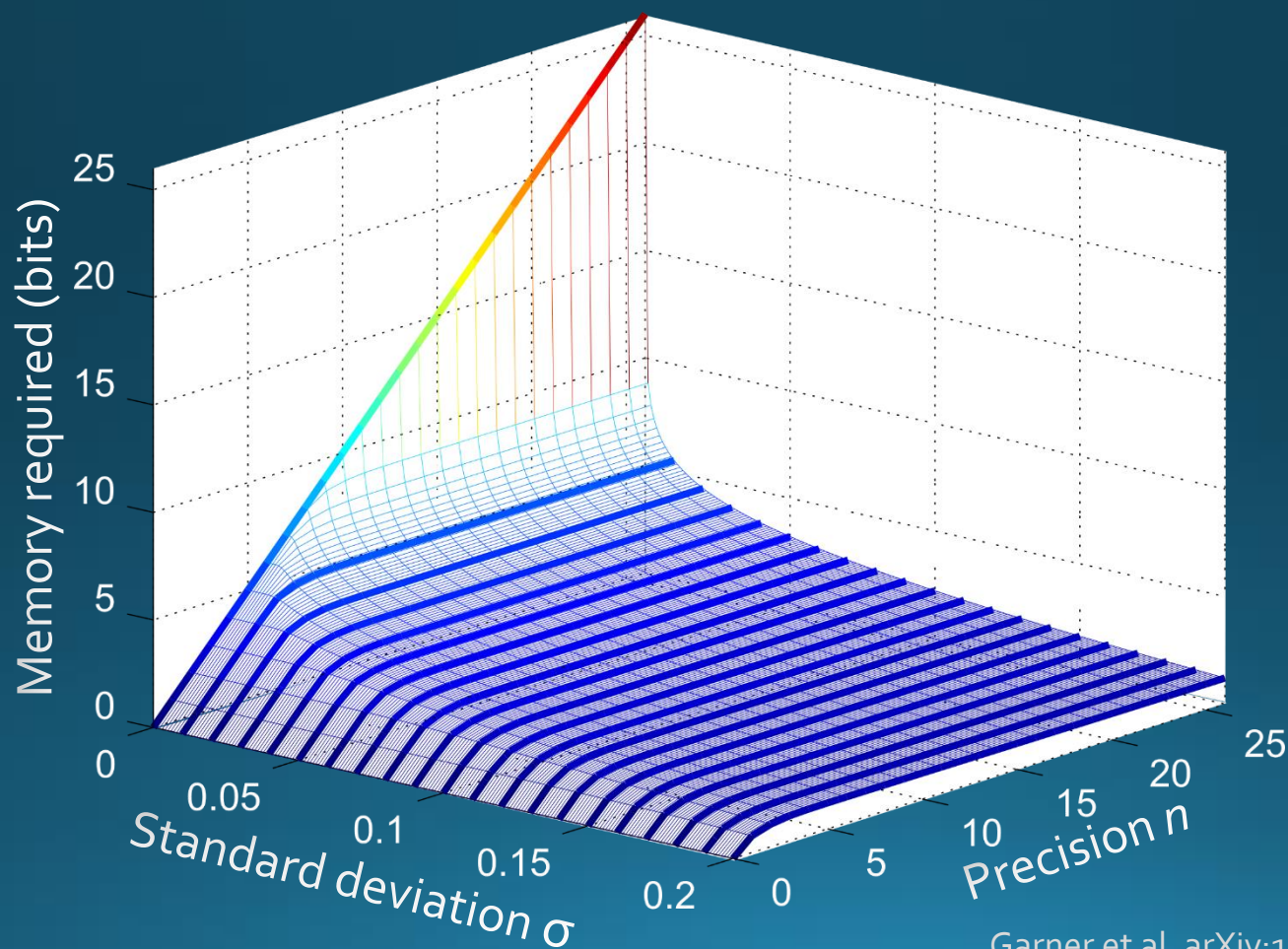
Shift function  $X$  describes Gaussian noise with standard deviation  $\sigma$ :

$$P(X = x) = \frac{1}{\sigma\sqrt{2\pi}} \exp\left(-\frac{(x - \mu)^2}{2\sigma^2}\right)$$



# Example: Gaussian noise

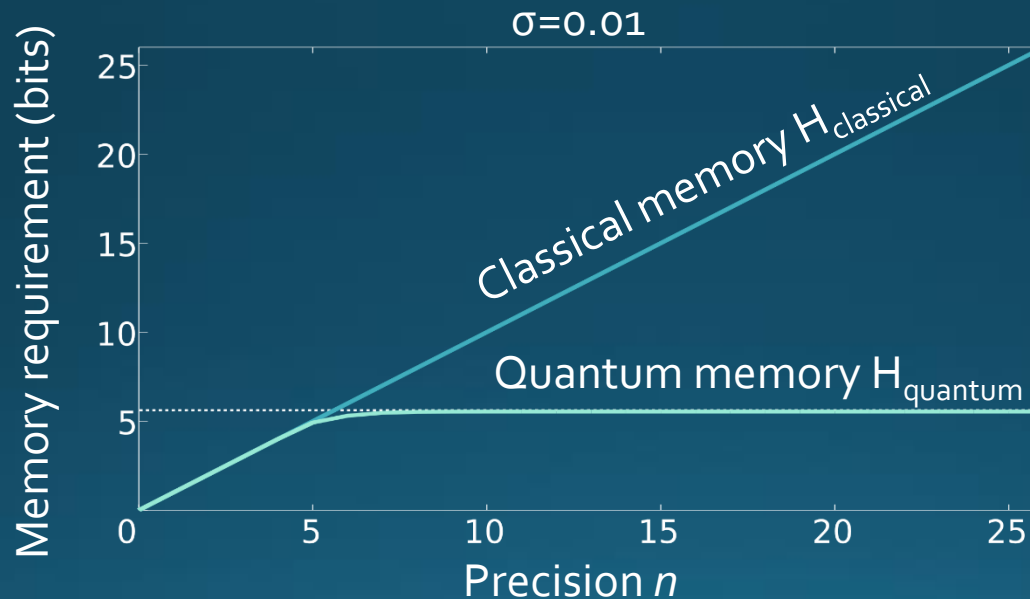
Entropy cost vs. precision:



# Example: Gaussian noise

Analytical bound for standard deviation  $0 < \sigma \ll 1$ :

$$\lim_{N \rightarrow \infty} H_{\text{quantum}} \leq \frac{1}{2 \ln 2} - (1 + 4\sqrt{2}\sigma) \log_2 \sigma$$



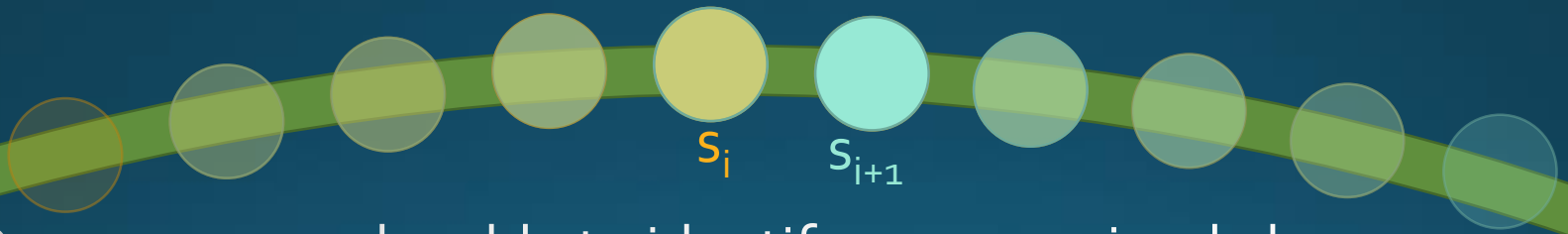
For small  $\sigma$ , leading term is  $-\log_2 \sigma$

# Origin of quantum advantage

As precision increases, the statistical future of neighbouring sites  $s_i, s_{i+1}$  becomes increasingly similar:

$$P(Y^{t+1}Y^{t+2} \dots | S^t = s_i) \approx P(Y^{t+1}Y^{t+2} \dots | S^t = s_{i+1})$$

but never actually converge for any finite precision  $n$ .



One may never be able to identify  $s_i$  vs.  $s_{i+1}$  simply by looking at future statistics  $Y^{t+1}Y^{t+2} \dots$  – the process exhibits unboundedly large **crypticity**.

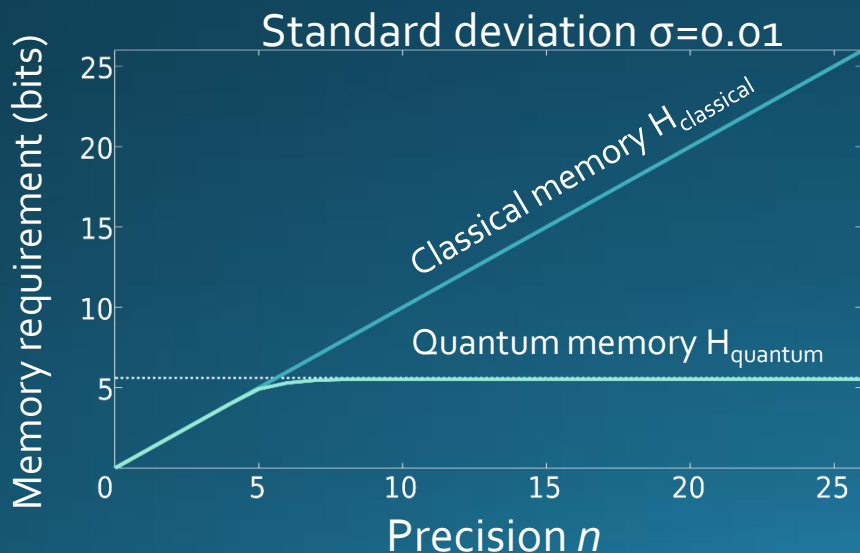
But, memory states  $s_i$  and  $s_{i+1}$  remain distinguishable! Hence, increasing precision **always** increases classical memory cost.

# Origin of quantum advantage

Increasing precision **always** increases classical memory cost.

As precision increases quantum states  $|S_i\rangle$  and  $|S_{i+1}\rangle$  increasingly overlap:  $\langle S_i | S_{i+1} \rangle \rightarrow 1$

Storing increasingly overlapping states adds a diminishing contribution to the memory's von Neumann entropy



Ultimately, precision can be increased with no further quantum memory cost.

# Observer-dependent complexity

The classical memory cost of simulation corresponds to its **statistical complexity**<sup>1</sup>.

Processes involving the stochastic dynamics of real numbers could be of very high complexity.

An observer with quantum information processing power could find such a process to be *much* simpler<sup>2</sup>.

**Quantum simulators** can simulate to arbitrarily high precision at fixed finite memory cost:

1. Crutchfield & Young, PRL (1989)

2. Suen et al. arXiv:1511.05738 (2015); Aghamohammadi et al. arXiv:1602.08646 (2016)

# Conclusions & outlook



# Conclusions

- Simulation of stochastic processes on classical digital computers requires a trade-off between precision and memory required.
- Quantum simulators can simulate to arbitrarily high precision at a *finite memory cost*.
- Quantum observers may find some processes simple, even if they exhibit unboundedly high classical complexity.

## Outlook

- These results probably generalize to any sort of simulation with *diffusive randomness*. (For example: non-Markovian or non-symmetric processes).
- This advantage may generalize to general *Monte Carlo* simulations, which could be run on a quantum processor to arbitrary precision, requiring a finite fixed memory cost.

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